

CPS 422
Computer Networks

Layered Approach To Network
Design

%% Networks are complex!

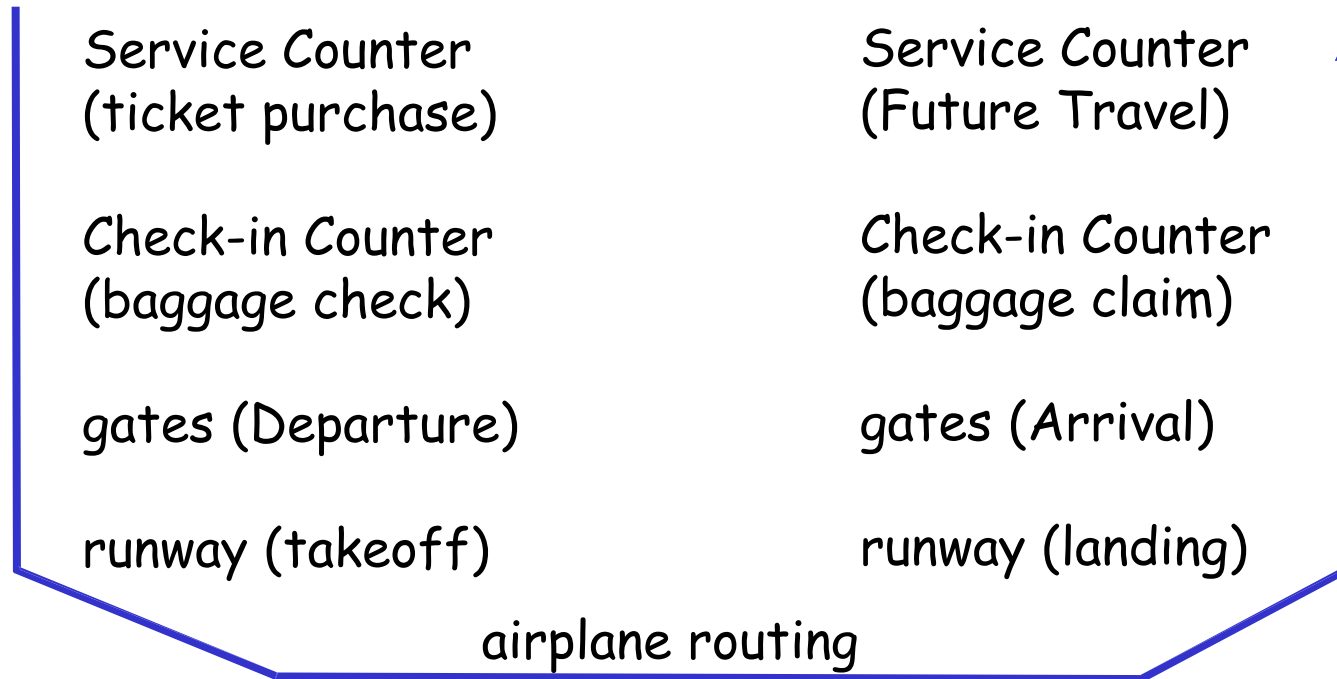
%% Composed of many "pieces":

- o hosts
- o routers
- o links of various media
- o applications
- o protocols
- o hardware, software

%% To better understand the issue let us first consider an analogy of "Air Travel Process"

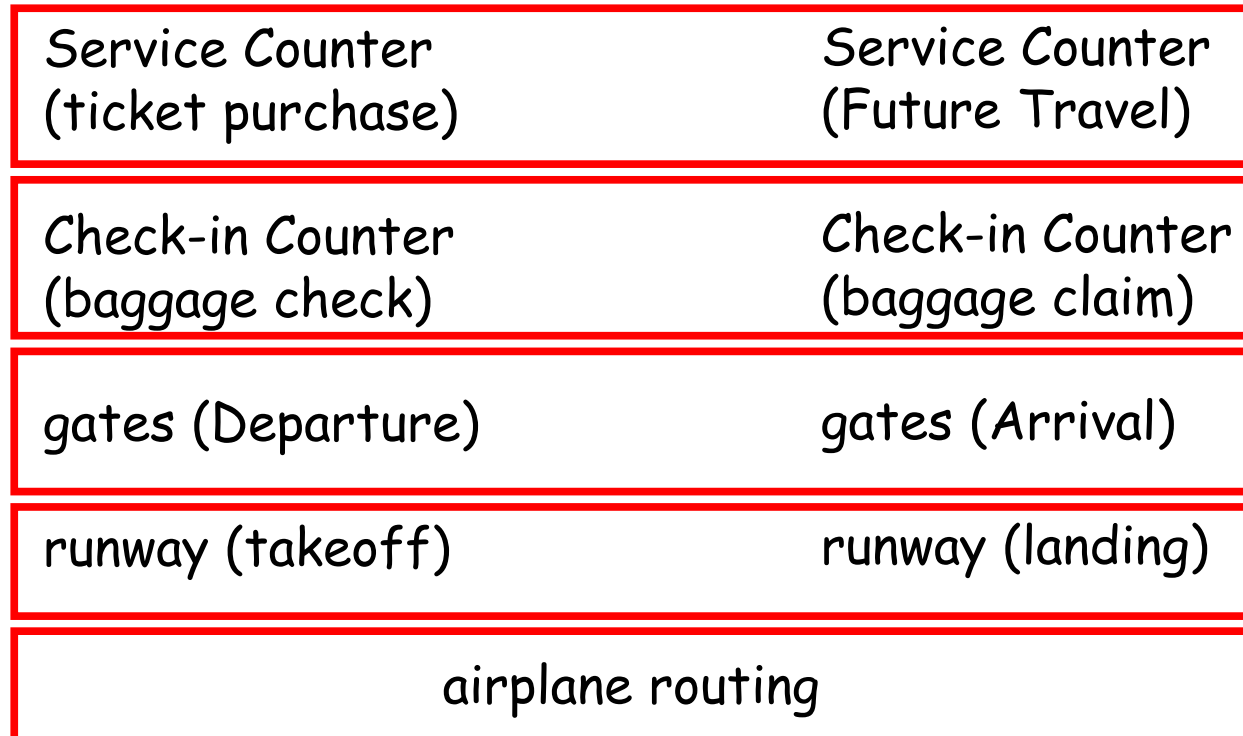
%% Can you break down the process of air travel into "sub-processes"

Organization of air travel



% a series of steps

Organization of air travel: a different view



Layers: each layer implements a service

- o via its own internal-layer actions
- o relying on services provided by layer below

Layered air travel: services

Counter-to-counter Ticketing services

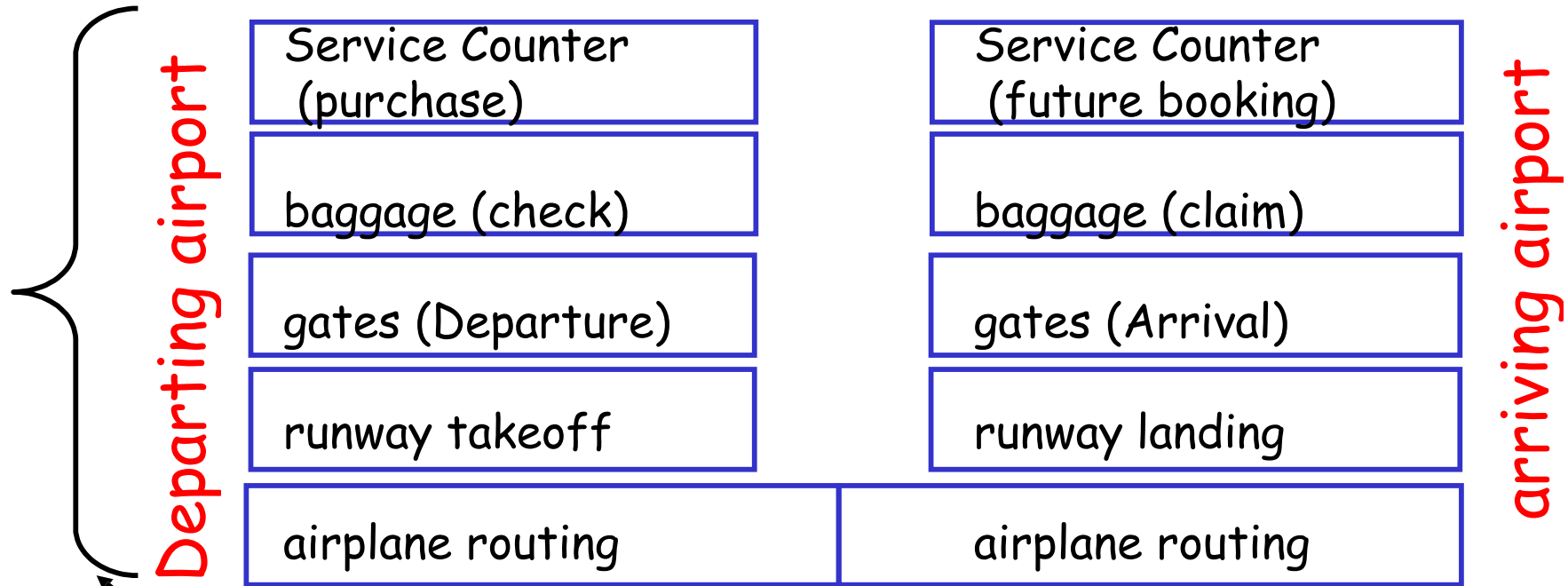
baggage-deposit-to-baggage-claim delivery

people transfer: Departure gate to Arrival gate

runway-to-runway delivery of plane

airplane routing from source to destination

Distributed implementation of layer functionality



We can refer to this collection as a "stack"

Similarly.....

‰ While developing a model for network architecture, designers distilled the process of transmitting data into its most fundamental elements.

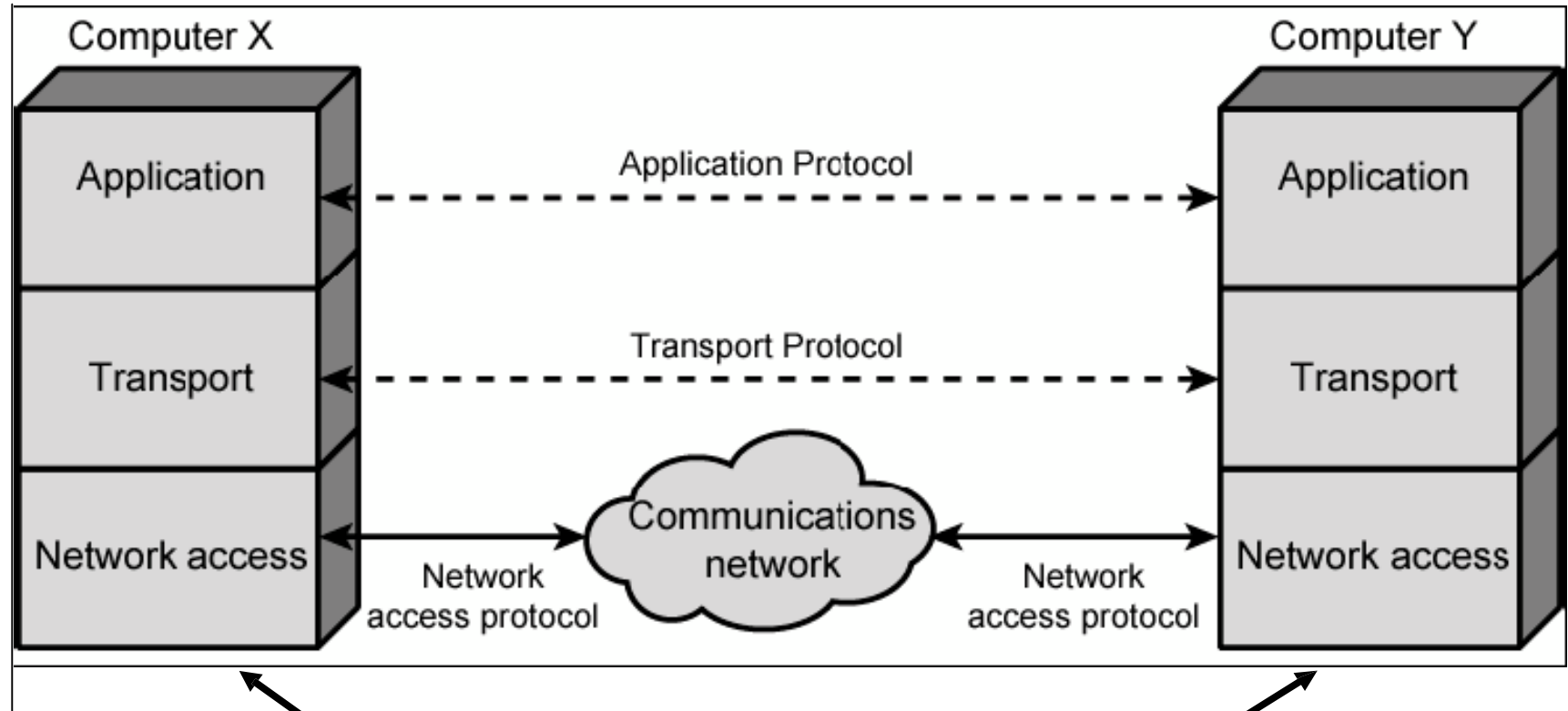
They identified which networking functions had related uses and collected those functions into discrete groups that became "Protocol Layers" which collectively form a "Protocol Stack".

Why layering?

Dealing with complex systems:

- ‰ explicit structure allows identification, relationship of complex system's pieces
 - layered **reference model** for discussion/design
- ‰ modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system

Protocols in Simplified Architecture

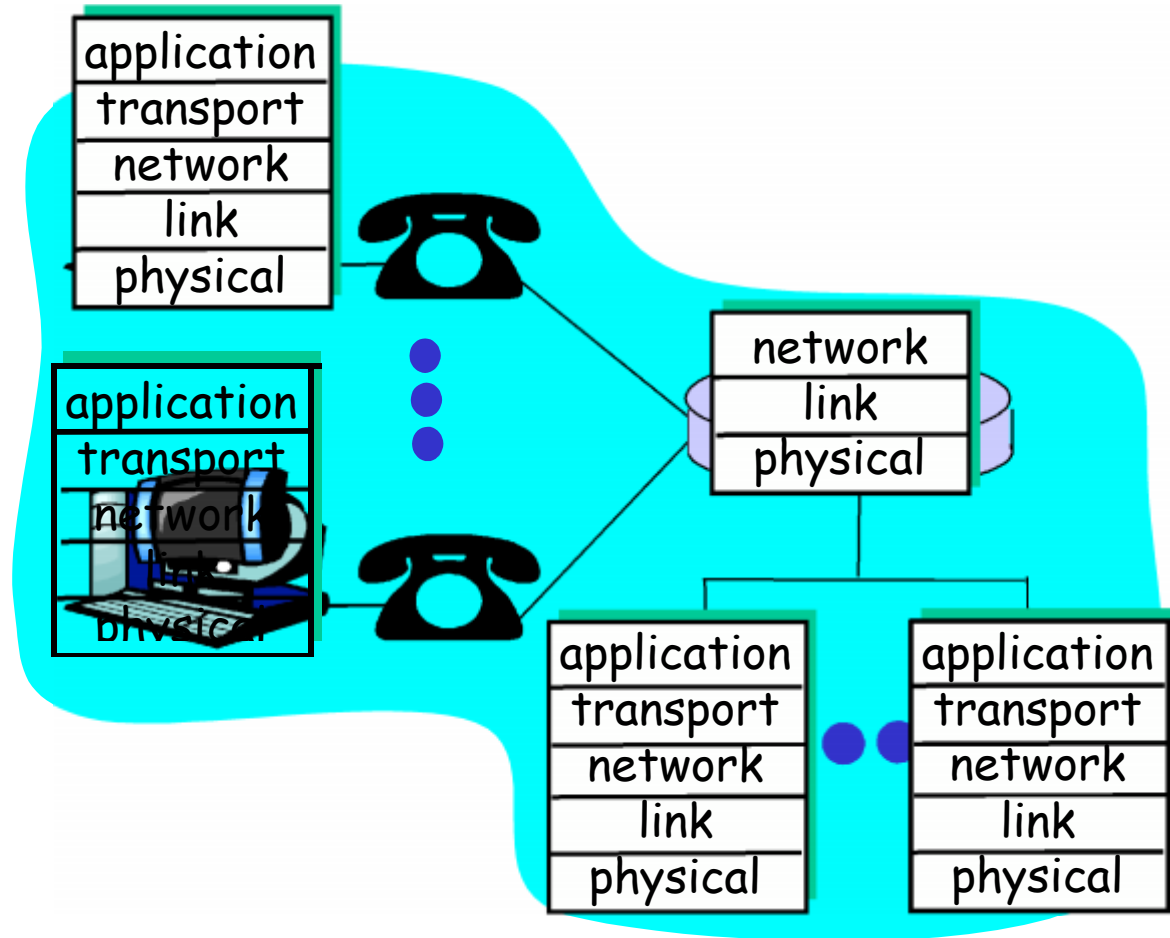


This is a Protocol Stack

Layering: logical communication

Each layer:

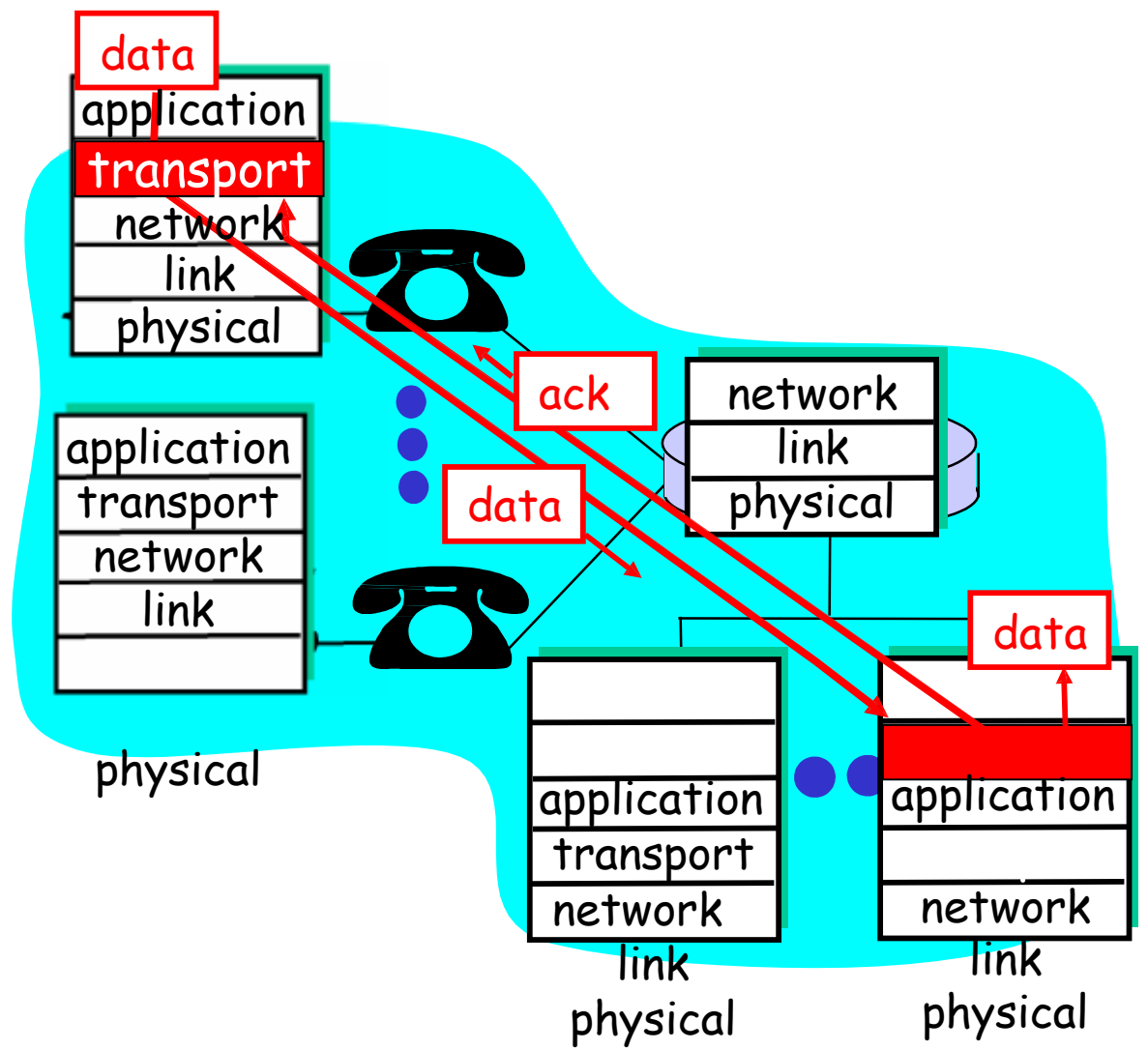
- distributed
- "entities" implement layer functions at each node
- entities perform actions, exchange messages with peers



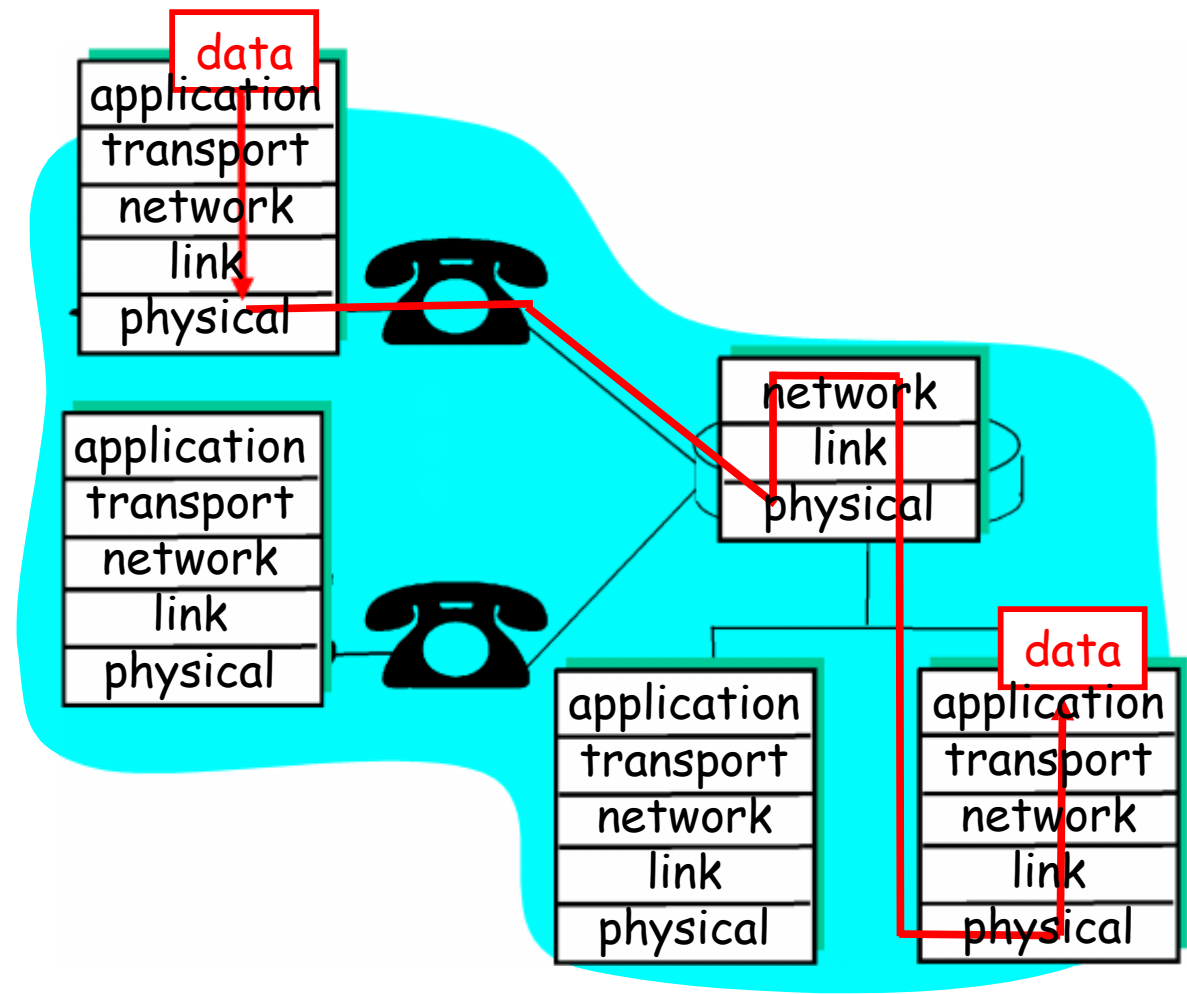
Layering: logical communication

E.g.: transport

- take data from app
- add addressing, reliability check info to form "datagram"
- send datagram to peer
- wait for ack receipt
- analogy: post office



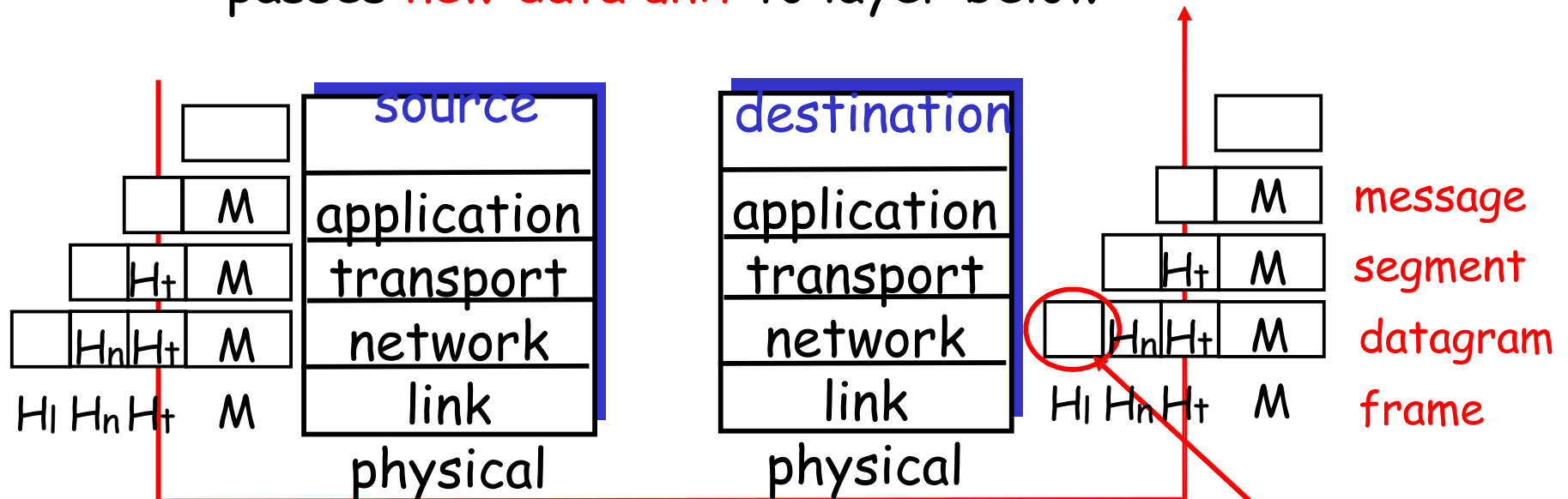
Layering: physical communication



Protocol layering and data

Each layer takes data from above

- ~ **adds header** information to create new data unit
- ~ passes **new data unit** to layer below



Protocol Revisited

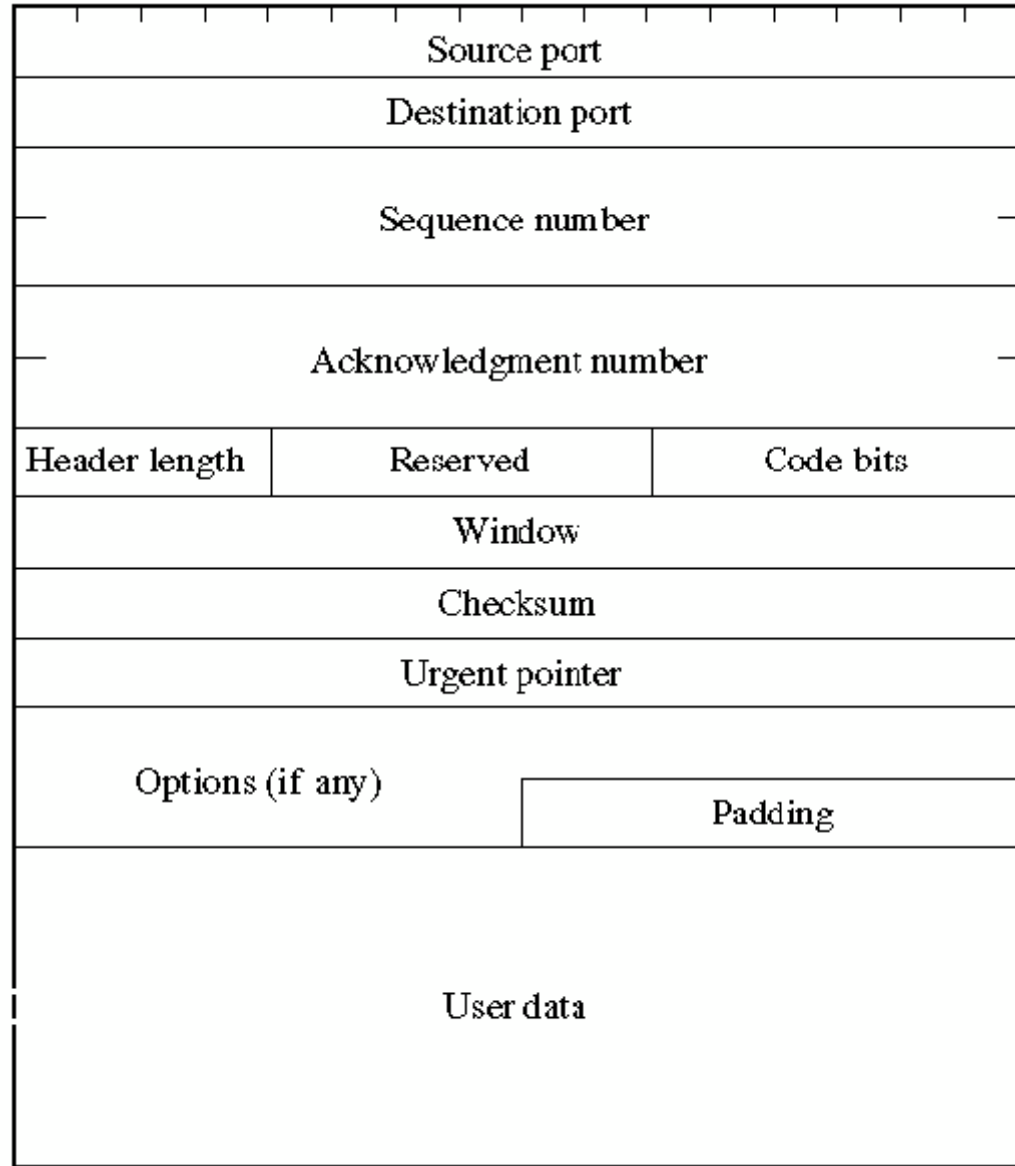
How does a Protocol Work

A Protocol implements its Functionality through
its **"HEADER"**

A Header contains all the control
information to complete the
protocol's tasks

TCP - Protocol Data Unit (PDU)

An - Example Header



%o Key Elements of a Protocol

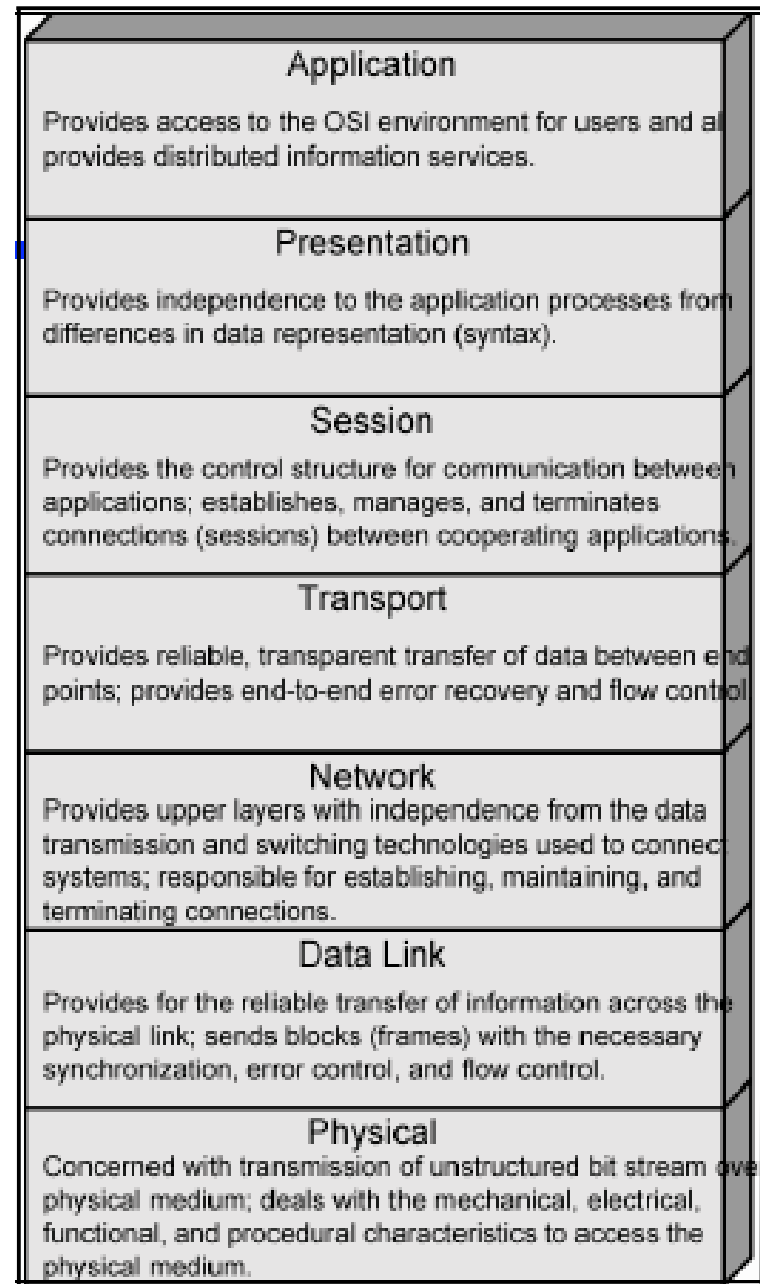
- o Syntax - Format
- o Semantics - Control Info for Coord and error handling
- o Timings - Sequencing and Speed matching

%o Categories of Protocol Functions

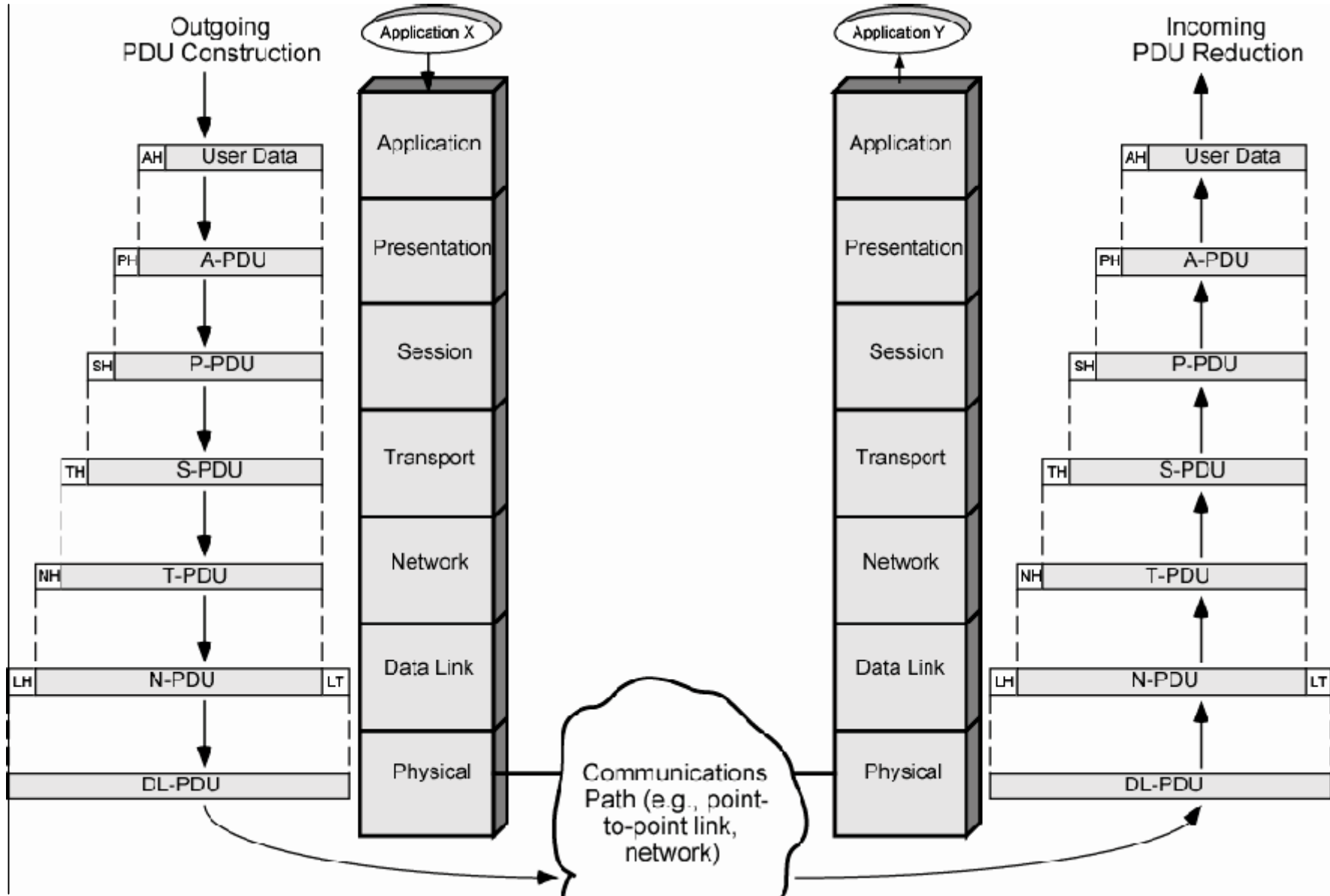
- o Segmentation and Re-assembly
- o Encapsulation
- o Connection Control
- o Ordered Delivery
- o Flow Control
- o Error Control
- o Addressing
- o Multiplexing
- o Transmission Services

The OSI Reference Model and TCP/IP Protocol Suite

The OSI Reference Model

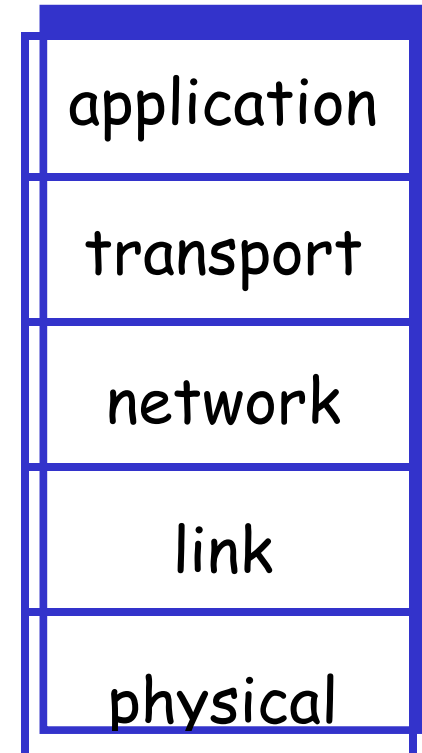


The OSI Environment

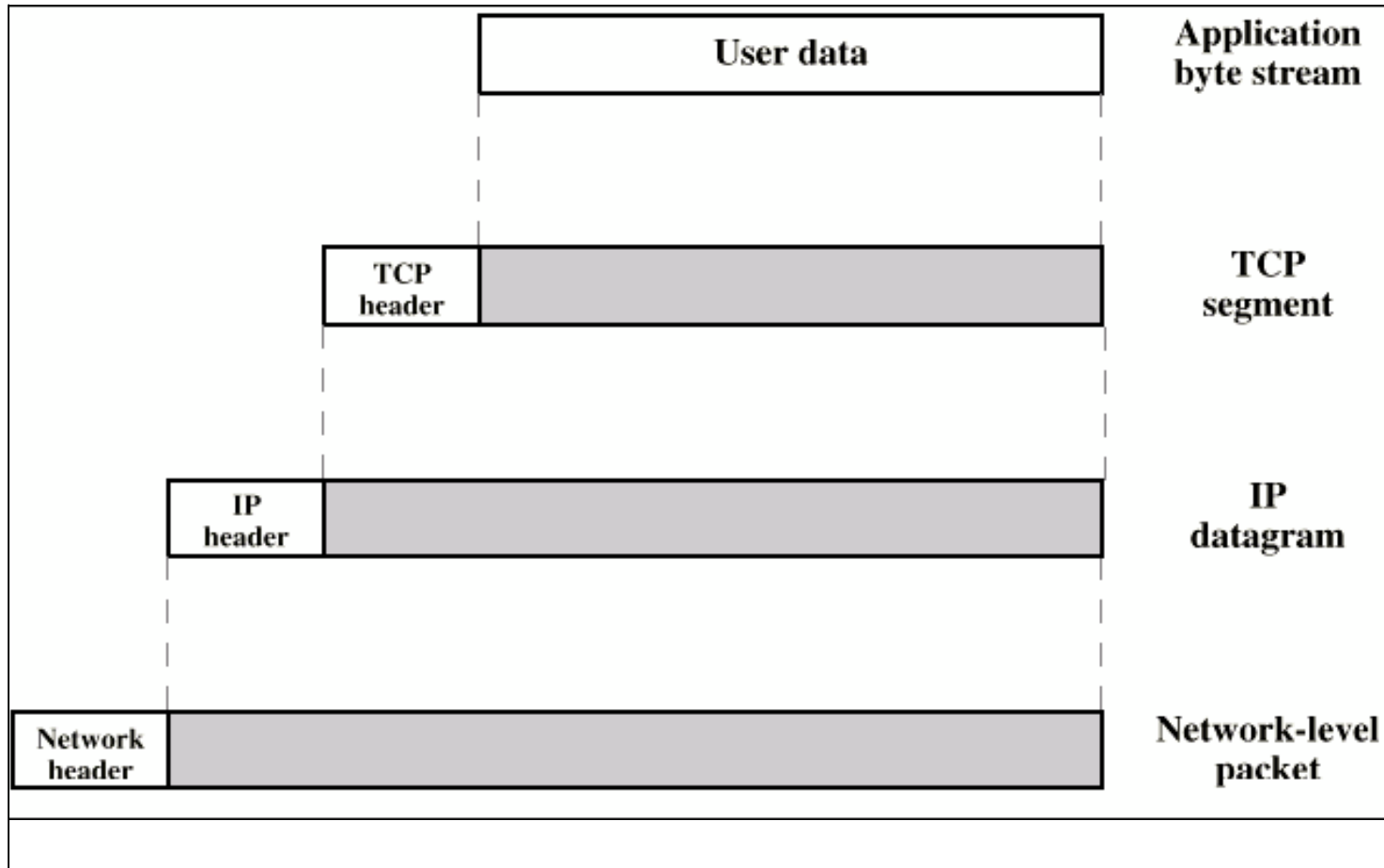


TCP / IP Protocol Suite

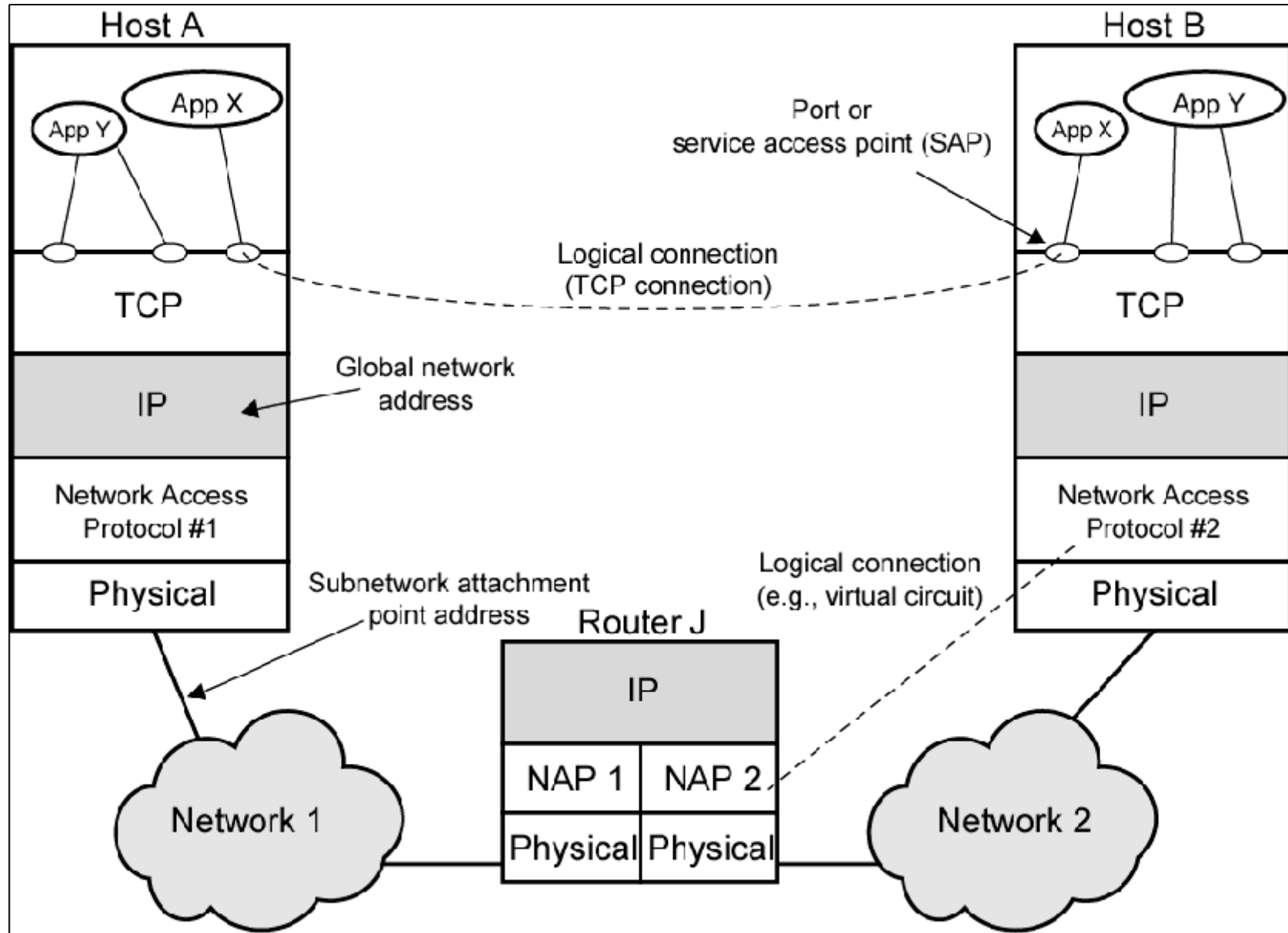
- ‰ **application:** supporting network applications
 - ftp, smtp, http
- ‰ **transport:** host-host data transfer
 - tcp, udp
- ‰ **network:** routing of datagrams from source to destination
 - ip, routing protocols
- ‰ **link:** data transfer between neighboring network elements
 - ppp, ethernet
- ‰ **physical:** bits "on the wire"



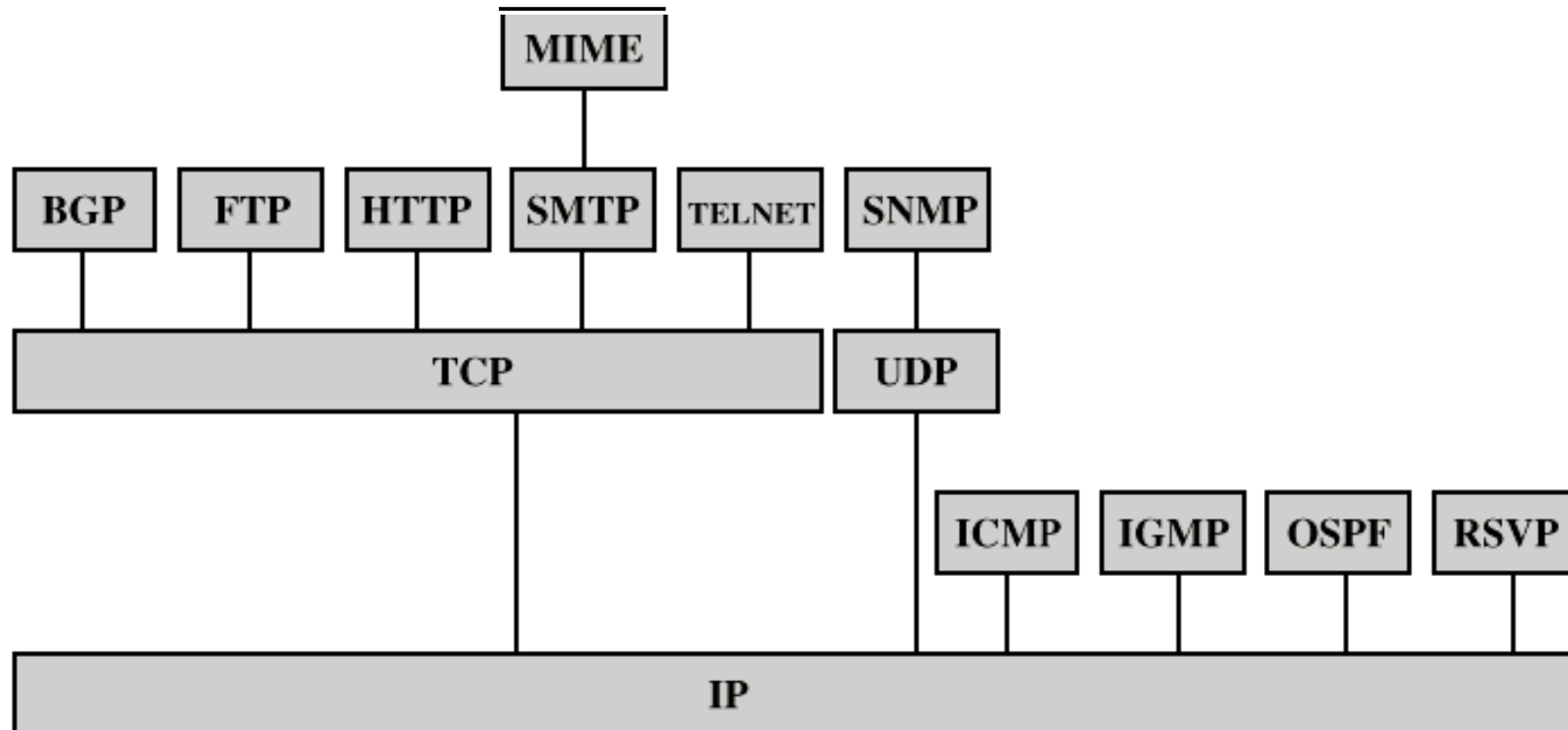
Some PDUs in TCP/IP



TCP / IP Concepts



Some Protocols of TCP/IP Suite



BGP = Border Gateway Protocol
FTP = File Transfer Protocol
HTTP = Hypertext Transfer Protocol
ICMP = Internet Control Message Protocol
IGMP = Internet Group Management Protocol
IP = Internet Protocol
MIME = Multi-Purpose Internet Mail Extension

OSPF = Open Shortest Path First
RSVP = Resource ReSerVation Protocol
SMTP = Simple Mail Transfer Protocol
SNMP = Simple Network Management Protocol
TCP = Transmission Control Protocol
UDP = User Datagram Protocol

OSI - TCP/IP Comparison

OSI	TCP/IP
Application	Application
Presentation	
Session	
Transport	Transport (host-to-host)
Network	Internet
Data Link	Network Access
Physical	Physical

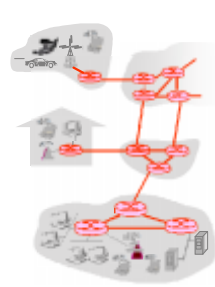
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Physical Layer Issues

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The Network Core

- ‰ mesh of interconnected routers
- ‰ **the fundamental question:** how is data transferred through the network?
- ™ **circuit switching:** dedicated circuit per call: telephone net
- ™ **packet-switching:** data sent through network in discrete "chunks"

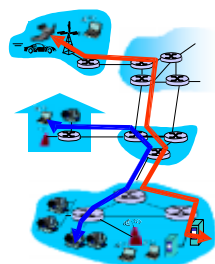


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Network Core: Circuit Switching

End-end resources reserved for "call"

- ‰ link bandwidth, switch capacity
- ‰ no sharing
- ‰ guaranteed performance
- ‰ call setup required



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Network Core: Circuit Switching

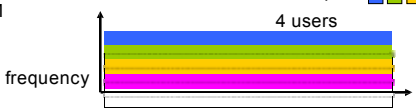
- ‰ network resources (e.g., bandwidth) **divided and allocated** to calls (Multiplexing)
- ‰ Resources remain **idle** if not used by owning call
- ‰ Techniques for dividing network resources into "pieces"
 - ™ frequency division
 - ™ time division
 - ™ Any other multiplexing techniques?

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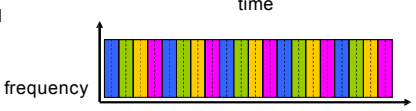
Circuit Switching: FDM and TDM

Example: ■ ■ ■ ■ 4 users

FDM



TDM



time

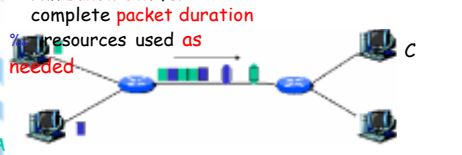
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Network Core: Packet Switching

each end-end data stream divided into packets

- ‰ packets from hosts share network resources
- ‰ each packet uses full link bandwidth for complete packet duration
- ‰ resources used as needed

Bandwidth division into "pieces"
Dedicated allocation Resource reservation



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Network Core: Packet Switching

resource contention:

- aggregate resource demand can exceed amount available
- congestion: packets queue-up, wait for link use
- store and forward: packets move one hop at a time

- Node receives complete packet before forwarding

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Packet switching versus circuit switching

Is packet switching a "clear winner?"

- Great for bursty data
- resource sharing
- simpler, no call setup
- Excessive congestion:** packet delay and loss
- protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior for a packet switched network?**
- bandwidth guarantees needed for audio/video apps
- QoS - Quality of Service (will be studied later in TCP)

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Packet-switched networks: Types

- Goal:** move packets from source to destination
- Two types of packet switch networks
- datagram network:**
 - destination address in packet determines next hop
 - routes may change during session
 - analogy: driving, asking directions
- virtual circuit network:**
 - each packet carries tag (virtual circuit ID), tag determines next hop
 - fixed path determined at call setup time, remains fixed throughout the session/call
 - Pro:** routers can perform resource reservation
 - Con:** routers maintain per-call state (complex)

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Delay, Loss and Throughput

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How do loss and delay occur?

packets queue-up in router buffers

- packet arrival rate to link exceeds output link capacity
- packets queue-up, wait for turn

packet being transmitted (delay)

packets queueing (delay)

free (available) buffers: arriving packets dropped (loss) if no free buffers

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Four sources of packet delay

- 1. nodal processing:**
 - check bit errors
 - determine output link
- 2. queueing:**
 - time spent waiting at output link for transmission
 - depends on congestion level of router

transmission

propagation

nodal processing

queueing

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Delay in packet-switched networks

3. Transmission delay:

- % R=link bandwidth (bps)
- % L=packet length (bits)
- % time to send bits into link = L/R

4. Propagation delay:

- % d = length of physical link
- % s = propagation speed in medium ($\sim 2-3 \times 10^8$ m/sec)
- % propagation delay = d/s

Note: s and R are very different quantities!

nodal processing queuing

Nodal delay

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

- % d_{proc} = processing delay
- % d_{queue} = queuing delay
- % d_{trans} = transmission delay
- % d_{prop} = propagation delay

"Real" Internet delays and routes

- % What do "real" Internet delay & loss look like?
- % **Traceroute program:** provides delay measurement from source to router along end-end Internet path towards destination. For all i:
 - TM sends three packets that will reach router i on path towards destination
 - TM router i will return packets to sender
 - TM sender times interval between transmission and reply.

Packet loss

- % queue (buffer) has finite capacity
- % when packet arrives to full queue, packet is dropped (lost)
- % Other sources of packet loss?
- % lost packet may be retransmitted by source (TCP), or not retransmitted at all (UDP)

Throughput

- % **throughput:** rate (bits/time unit) at which bits transferred between sender/receiver
- TM **instantaneous:** rate at given point in time
- TM **average:** rate over long(er) period of time

server sends bits (fluid) into pipe pipe that can carry bits (fluid) at rate (R_s bits/sec) pipe that can carry (bits) fluid at rate (R_c bits/sec)

Throughput

What is end-end throughput?

R_s bits/sec R_c bits/sec

What is end-end throughput?

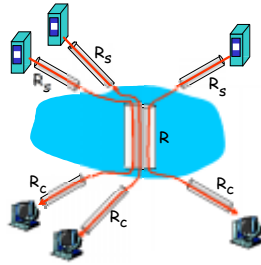
R_s bits/sec R_c bits/sec

bottleneck link

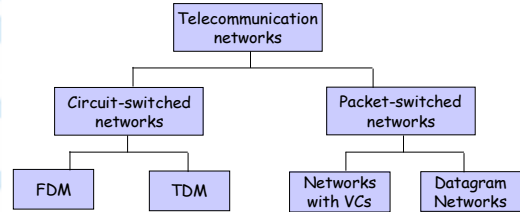
link on end-end path that constrains end-end throughput

Throughput: Internet scenario

- per-connection end-end throughput: $\min(R_c, R_s)$
- R_c or R_s often the bottleneck link
- Is it the case nowadays?



Network Taxonomy

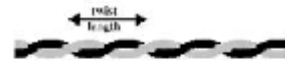


- Internet is a datagram / packet-switched network
- Internet provides both **connection-oriented** (TCP) and **connectionless services** (UDP) to applications.

Transmission Media

Twisted Pair (TP)

- Separately insulated
- Twisted together
- Often "bundled" into cables
- Usually installed in building during construction



(a) Twisted pair

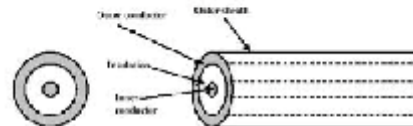
- Twists decrease the cross-talk
- Neighboring pairs have different twist length
- Most of telephone and network wiring in homes and offices is TP.



Unshielded and Shielded TP

- Unshielded Twisted Pair (UTP)
 - Ordinary telephone wire
 - Cheap, Flexible \Rightarrow Easiest to install
 - No shielding \Rightarrow Suffers from external EM interference
 - Used in Telephone and Ethernet
- Shielded Twisted Pair (STP)
 - Metal braid or sheathing that reduces interference
 - More expensive
 - Harder to handle (thick, heavy)
 - Used in token rings

Coaxial Cable



- Outer conductor is braided shield
- Have used in cable TV case
- Insulated by insulating material
- Covered by plastic

- Higher bandwidth than UTP. Up to 500 MHz.
- Used in cable TV

Reflection and Refraction

- ❑ Index of Refraction = Speed of light in Vacuum/Speed in glass
= $300 \text{ m}/\mu\text{s} / 200 \text{ m}/\mu\text{s} = 1.5$

Index = 1.48 Refraction

Index = 1.5 Reflection

- ❑ Refracted light bends towards the higher index medium

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Optical Fiber

- ❑ A cylindrical mirror is formed by the cladding
- ❑ The light wave propagate by continuous reflection in the fiber
- ❑ Not affected by external interference \Rightarrow low bit error rate
- ❑ Fiber is used in all long-haul or high-speed communication

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Types of Fibers

- ❑ Multimode Fiber: Core Diameter 50 or 62.5 μm
Wide core \Rightarrow Several rays (mode) enter the fiber
Each mode travels a different distance
- ❑ Single Mode Fiber: 10- μm core. Lower dispersion.

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Antenna

- ❑ Transmitter converts electrical energy to electromagnetic waves
- ❑ Receiver converts electromagnetic waves to electrical energy
- ❑ Same antenna is used for transmission and reception
- ❑ Omni-Directional: Power radiated in all directions
- ❑ Directional: Most power in the desired direction
- ❑ Isotropic antenna: Radiates in all directions equally
- ❑ Antenna Gain = Power at particular point/Power with Isotropic Expressed in dBi

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Terrestrial Microwave

- ❑ Parabolic dish
- ❑ Focused beam
- ❑ Line of sight
- ❑ Long haul telecommunications
- ❑ Higher frequencies give higher data rates

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Satellite Microwave

- ❑ Relay station \Rightarrow Satellite receives on one frequency, amplifies or repeats signal and transmits on another frequency
- ❑ Geostationary orbit: Height of 35,784 km
- ❑ Point to Point or Direct broadcast satellite

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Infrared

- ❑ Used in TV remote control
- ❑ IRD port of computers
- ❑ Modulate infrared light
- ❑ Line of sight (or reflection)
- ❑ Blocked by walls



Wireless Propagation

- ❑ **Ground wave:** Follows contour of earth. Up to 2MHz. AM radio
- ❑ **Sky wave:** Signal reflected (Actually refracted) from ionosphere layer of upper atmosphere.
- ❑ **Line of sight:** Above 30MHz.

